The Ghost Algebra

Madeline Nurcombe

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12 May 2023

1 Why do we want a new algebra?

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- 2 What is the ghost algebra?
- 3 What is the dilute ghost algebra?
- **4** What can we do with these algebras?

QUESTION 1:

Why do we want a new algebra?

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and **unital**, meaning there exists $I \in A$ such that

$$la = al = a, \quad \forall a \in A.$$

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Equality up to isotopy:





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- Dimensions given by Catalan numbers

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• Associativity implies $\alpha_1 = \alpha_2$

QUESTION 2: What is the ghost algebra?

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The Ghost Algebra

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Equality: string isotopy; same # ghosts in each boundary region mod 2















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Evenness condition in TL_n^2

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What if some nodes didn't have any strings connected to them?

QUESTION 3: What is the dilute ghost algebra?

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The Ghost Algebra

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$$\left(\begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix} \right) \begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 2 & 3 \\ 0 & 1 \end{bmatrix} = \alpha_1 \begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix}$$

Boundary associativity problem occurs even in one-boundary case.

What about one- or two-boundary dilute TL algebras?

$$\left(\begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix} \right) \begin{bmatrix} 1 & 2 & 3 & 4 & 5 \\ 0 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2 \begin{bmatrix} 1 & 2 & 3 & 4 \\ 0 & 1 & 1 \end{bmatrix} = \alpha_2$$

Boundary associativity problem occurs even in one-boundary case.

No algebras in literature have $dTL_n(\beta)$ and $TL_n^1(\beta; \alpha_1, \alpha_2)$ as subalgebras with $\alpha_1 \neq \alpha_2$.

Dilute ghost algebra dGh_n^2 allows *n*-diagrams with empty nodes

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String to empty node gives 0:

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$$\operatorname{Gh}^2_{n} \begin{pmatrix} \alpha_1 & \alpha_2 & \alpha_3 \\ \beta & \gamma_{12} & \gamma_3 \\ \delta_1 & \delta_2 & \delta_3 \end{pmatrix}$$

$$\operatorname{Gh}^{1}_{n}(\beta; \alpha_{1}, \alpha_{2}, \alpha_{3}) \longleftrightarrow \operatorname{Gh}^{2}_{n} \begin{pmatrix} \alpha_{1} & \alpha_{2} & \alpha_{3} \\ \beta & \gamma_{12} & \gamma_{3} \\ \delta_{1} & \delta_{2} & \delta_{3} \end{pmatrix}$$

$$\operatorname{TL}_{n}^{1}(\beta;\alpha_{1},\alpha_{2})$$

$$\int$$

$$\operatorname{Gh}_{n}^{1}(\beta;\alpha_{1},\alpha_{2},\alpha_{3}) \longleftrightarrow \operatorname{Gh}_{n}^{2}\begin{pmatrix}\alpha_{1} & \alpha_{2} & \alpha_{3} \\ \beta & \gamma_{12} & \gamma_{3} \\ \delta_{1} & \delta_{2} & \delta_{3} \end{pmatrix}$$

$$TL_{n}(\beta) \longleftrightarrow TL_{n}^{1}(\beta; \alpha_{1}, \alpha_{2})$$

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QUESTION 4:

What can we do with these algebras?

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Build a lattice out of **bulk squares** and **boundary triangles**



Build a lattice out of bulk squares and boundary triangles



Construct Hamiltonians using solutions to **Yang-Baxter equation** (YBE) and **boundary Yang-Baxter equation** (BYBE).

Build a lattice out of bulk squares and boundary triangles



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Construct Hamiltonians using solutions to Yang-Baxter equation (YBE) and boundary Yang-Baxter equation (BYBE). YBE:

$$u = \frac{u}{v} = \frac{u}{v} + \frac{\sin(u)}{\sin(\lambda)} + \frac{\sin(u)}{\sin(\lambda)} + \beta = 2\cos\lambda$$

Boundary Yang-Baxter equation

The BYBE involves boundary face operators

$$u - v \qquad u + v \qquad = \qquad v \qquad u + v \qquad u - v$$

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Two solutions: $\rho \in \mathbb{C}$, $\beta = 2 \cos \lambda$, $b : \mathbb{C} \to \mathbb{C}$,

$$u = b(u) \left(\frac{F(u)}{\sin(2u)} + c_1 + c_2 + c_3 + c_4 +$$

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$$\begin{aligned} F(u) &= \rho + \left(2\alpha_3^2 - \alpha_1^2 - \alpha_2^2\right)\cos(2u) - \left(\alpha_3^2 - \alpha_1\alpha_2\right)\cos(2u - \lambda), \\ c_1 &= -\alpha_1, \qquad c_2 = c_4 = \alpha_3, \qquad c_3 = -\alpha_2. \end{aligned}$$

Dilute loop model



Dilute loop model



Batchelor & Yung (1995) — YBE solution ($\beta = -2\cos(4\phi)$)

$$u := w_1(u) + w_2(u) \left(+ \cdots \right) + w_3(u) \left(+ \cdots \right) + w_4(u) \left(+ \cdots \right) + w_5(u) + w_6(u) \cdots + w_6(u) \right)$$

$$\begin{split} w_1(u) &= \sin(2\phi)\sin(3\phi) + \sin u \sin(3\phi - u) & w_2(u) = \sin(2\phi)\sin(3\phi - u) \\ w_3(u) &= \sin(2\phi)\sin u & w_4(u) = \sin u \sin(3\phi - u) \\ w_5(u) &= \sin(2\phi - u)\sin(3\phi - u) & w_6(u) = -\sin u \sin(\phi - u) \end{split}$$

Dilute ghost algebra BYBE solution

$$u = a(u) + b_1(u) + b_2(u) + b_3(u) + b_4(u) + b_5(u) + b_6(u) + b_7(u) + b_8(u) + b_9(u) +$$

Dilute ghost algebra BYBE solution

$$u = a(u) + b_1(u) + b_2(u) + b_3(u) + b_4(u) + b_4(u) + b_5(u) + b_6(u) + b_7(u) + b_8(u) + b_9(u) +$$

$$\begin{split} & \mathbf{d}_{0} = -\frac{\cos(\alpha_{0} \left(x-\frac{\alpha_{0}}{2}\right)}{\cos(\alpha_{0} \left(x+\frac{\alpha_{0}}{2}\right)} \left(z_{0}^{*} \operatorname{cn}(\alpha) - x^{*} \left(\left(x^{*} + x^{*}\right) \alpha_{0} + \mu_{0} \varepsilon(\alpha_{0} + \alpha_{0})\right) \operatorname{dn}^{*} \left(x-\frac{\alpha_{0}}{2}\right)\right) \\ & \mathbf{h}(\alpha) = -\frac{\mathbf{d}_{0}}{\sin(\alpha_{0} \cos(\alpha_{0})} \left(z^{*} \operatorname{cn}^{*}\left(x+\frac{\alpha_{0}}{2}\right) \operatorname{cn}(\alpha) - x^{*} \left(\left(x^{*} + x^{*}\right) \alpha_{0} + \mu_{0} \varepsilon(\alpha_{0} + \alpha_{0})\right) \operatorname{cn} \left(x - \frac{\alpha_{0}}{2}\right) \operatorname{cn} \left(x - \frac{\alpha_{0}}{2}\right) \right) \\ & \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = -\operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = -\operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = -\operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = -\operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \mathbf{h}(\alpha) = -\operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \qquad \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \quad \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \quad \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \quad \mathbf{h}(\alpha) = \operatorname{cn}^{*}(\alpha), \\ & \mathbf{h}(\alpha) = \operatorname{cn}^{*}($$

$$\begin{split} d_{2}(q) &= \frac{2\pi q}{m_{2}(q)}\left(m_{1} - m_{2}\right) \left(m_{1}(m_{1} + m_{2}) + \left(n^{2} + n^{2}\right) m_{2}\right) \left(m_{2}(m_{2}(q) - m_{2}(q)) \left(m_{2}(m_{2} - q) - m_{2}(q) - m_{2}(q) \right) \right) \\ h_{1}(q) &= \frac{2\pi q}{m_{2}(q)}\left(m_{1} - m_{2}\right) \left(m_{1}(m_{1} + m_{2}) + \left(n^{2} + n^{2}\right) m_{2}\right) \left(m_{2}(m_{2}(q) + m_{2}) \right) \left(m_{2}(m_{2} - q) - m_{2}(q) - m_{2}(q) \right) \right) \\ h_{1}(q) &= h_{1}(q) = h_{1}(q) = h_{1}(q) = h_{1}(q) = h_{1}(q) \\ h_{2}(q) &= h_{2}(q) \left(m_{1}(m_{2} + m_{2}) \left(m_{2}(m_{2}(q) - m_{2}) \left(m_{2}(m_{2}(m_{2}(q) - m_{2}) \left(m_{2}(m_{2}(m_{2}(q) - m_{2}) \left(m_{2}(m_$$

$$\begin{split} &g(z) = \frac{(z-z)}{4\pi (2q-4z)} \left[2q^2 \sin(z) - a^2 \left(\left(a^2 + a^2 \right) z_0 - a u(z_0 + z_0) \right) m^2 \left(a - \frac{2}{2} \right) \right) \\ &h(z) = \frac{-6q}{4\pi (2q-4z)} \left(2q^2 m \left(c + \frac{2q}{2} \right) m(z) - a^2 \left(\left(a^2 + a^2 \right) z_0 - a u(z_0 + z_0) \right) m \left(c + \frac{2}{2} \right) m \left(c - \frac{2}{2} \right) \right) \\ &h(z) = e^{-aq^2} \left(b_1 \right) = -e^{-aq^2} \left(b_1 \right) = -e^{-aq^2} \left(b_1 \right) = e^{-aq^2} \left(b_1 \right) = e^{-aq^2} \left(b_1 \right) = -e^{-aq^2} \left(b_1 \right) = -e^{-aq^2} \left(b_1 \right) = -e^{-aq^2} \left(b_1 \right) = e^{-aq^2} \left(b_1 \right$$

$$\begin{split} d_{0}(s) &= -\frac{n_{0}(s)}{2\pi\omega(2s)} - \left(\frac{s^{2}}{2}, \frac{s^{2}}{2}, \frac{$$

 $\left(k^2-1\right)\left(\left(\alpha_1\,\alpha_2\,-\,\alpha_3^2\right)\sin(3\phi)-\left(\alpha_1^2+\alpha_2^2-2\alpha_3^2\right)\sin(\phi)\right)=2k\left(\left(\alpha_1\,\alpha_2\,-\,\alpha_3^2\right)\cos(3\phi)+\left(\alpha_1^2+\alpha_2^2-2\alpha_3^2\right)\cos(\phi)\right)$

$$\begin{split} g_{(1)} &= \frac{(a_{1},a_{1},a_{1},a_{2},a_{3}) + (a_{2}-a_{1})(a_{2}-a_{3})(a_{2$$

Madeline Nurcombe (UQ)

The Ghost Algebra

• Representation theory — use cellularity

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- Physics Hamiltonians and energy eigenvalues

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Ghost algebra generators



Dimensions

$$\dim \operatorname{Gh}_{n}^{1} = \sum_{d=0}^{n} \left(\sum_{j=0}^{\lfloor \frac{n-d}{2} \rfloor} 2^{n-2j-d} \left(\binom{n}{j} - \binom{n}{j-1} \right) \right)^{2}$$
$$\dim \operatorname{Gh}_{n}^{2} = \sum_{d=0}^{n} \left(\sum_{j=0}^{\lfloor \frac{n-d}{2} \rfloor} 2^{n-2j-d} (n-2j-d+1) \left(\binom{n}{j} - \binom{n}{j-1} \right) \right)^{2}$$
$$\dim \operatorname{Gh}_{n}^{1} = \sum_{d=0}^{n} \left(\sum_{v=0}^{n-d} \binom{n}{v} \sum_{j=0}^{\lfloor \frac{n-v-d}{2} \rfloor} 2^{n-v-2j-d} \left(\binom{n-v}{j} - \binom{n-v}{j-1} \right) \right)^{2}$$
$$\dim \operatorname{Gh}_{n}^{2} = \sum_{d=0}^{n} \left(\sum_{v=0}^{n-d} \binom{n}{v} \sum_{j=0}^{\lfloor \frac{n-v-d}{2} \rfloor} 2^{n-v-2j-d} (n-v-2j-d+1) \times \left(\binom{n-v}{j} - \binom{n-v}{j-1} \right) \right)^{2}$$

n	dim Gh_n^1	dim Gh_n^2	dim dGh ¹ _n	dim dGh_n^2
1	5	17	10	26
2	30	186	117	521
3	185	1,813	1,407	9,355
4	1,150	16,102	17,083	156,947
5	7,170	135,866	208,284	2,514,932
6	44,760	1,099,276	2,544,751	38,968,815
7	279,585	8,639,133	31,125,138	588,475,298
8	1,746,870	66,258,526	380,928,795	8,706,799,523
9	10,916,150	498,701,470	4,663,705,782	126,690,947,758
10	68,219,860	3,693,607,300	57,109,857,519	1,818,028,127,339