An Unexpected Equivalence From A Silly Cross Product Puzzle

Max Orchard

August 26, 2022

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An Unexpected Equivalence

August 26, 2022 1 / 21

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Inspiration

A post on the r/math subreddit:

What are some of the most interesting equivalent statements in math that intrigue(d) you? submitted 28 days ago by Colver_4k Algebra \mathfrak{B}

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A paper by Kauffman:

JOURNAL OF COMBINATORIAL THEORY, Series B 48, 145-154 (1990)

[SPOILER] and the Vector Cross Product*

LOUIS H. KAUFFMAN

Department of Mathematics, Statistics, and Computer Science, University of Illinois at Chicago, Chicago, Illinois 60680

Communicated by the Managing Editors

Received July 14, 1987; revised July 4, 1988

[SPOILER] is equivalent to a combinatorial problem about the three-dimensional vector cross product algebra. © 1990 Academic Press, Inc.

An Unexpected Equivalence

Let $\{i, j, k\}$ denote the standard basis for \mathbb{R}^3 .

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The cross product is defined on this basis by

$$i \times i = j \times j = k \times k = 0,$$

$$i \times j = k, \quad j \times k = i, \quad k \times i = j,$$

$$j \times i = -k, \quad k \times j = -i, \quad i \times k = -j,$$

and extended to $\mathbb{R}^3\times\mathbb{R}^3$ using bilinearity.

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We will only be looking at the cross product on $\{\pm i, \pm j, \pm k, 0\}$.

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The Puzzle

The cross product is not associative. For example,

$$(i \times i) \times j = 0 \times j = 0, \quad i \times (i \times j) = i \times k = -j.$$

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Off-topic remark: the number of associations of n variables is equal to the n^{th} Catalan number. The sequence of Catalan numbers is

 $1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, \ldots$

They grow fairly quickly!

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Fix two associations L and R of variables X_1, \ldots, X_n .

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If we choose each X_i to be a vector from the set $\{i, j, k\}$, the result of $L(X_1, \ldots, X_n)$ and $R(X_1, \ldots, X_n)$ will lie in $\{\pm i, \pm j, \pm k, 0\}$.

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Goal

Find solutions to the equation $L(X_1, \ldots, X_n) = R(X_1, \ldots, X_n)$.

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Find *sharp* solutions to the equation $L(X_1, \ldots, X_n) = R(X_1, \ldots, X_n)$.

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The Equivalence

We can always find a sharp solution for n = 3.

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We can always find a sharp solution for n = 3. There are only two distinct associations, given by

 $L(X_1, X_2, X_3) = (X_1 \times X_2) \times X_3, \qquad R(X_1, X_2, X_3) = X_1 \times (X_2 \times X_3).$

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Theorem (Kauffman)

The existence of a sharp solution to the equation L = R for any $n \in \mathbb{Z}^+$ and for all associations L, R of the variables X_1, \ldots, X_n is equivalent to

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The existence of a sharp solution to the equation L = R for any $n \in \mathbb{Z}^+$ and for all associations L, R of the variables X_1, \ldots, X_n is equivalent to the four colour theorem.

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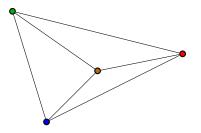
Theorem (Four Colour Theorem)

Every simple planar graph can be vertex-coloured with four colours.

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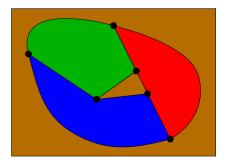
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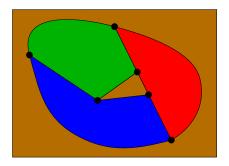


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Theorem (Four Colour Theorem)

Every bridgeless cubic planar graph can be face-coloured with four colours.



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We can reformulate the four colour theorem into an edge-colouring problem.

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Theorem

Every bridgeless cubic planar graph can be edge-coloured with three colours.

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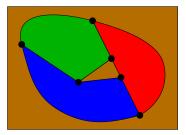
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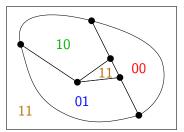
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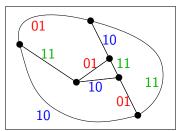
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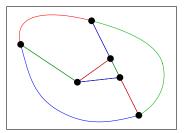
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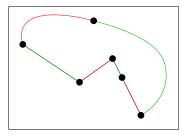
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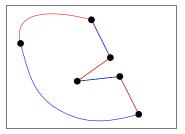
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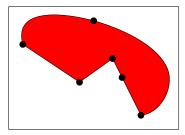


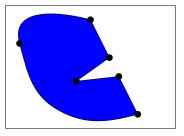
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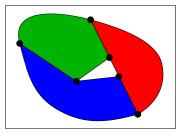


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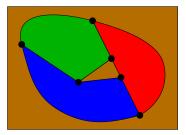
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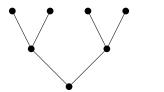
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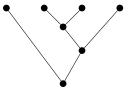
Let L and R be two associations of X_1, \ldots, X_n . We can construct a tree from an association by pairing up each individual multiplication.

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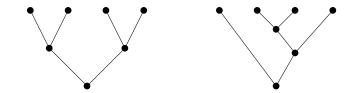
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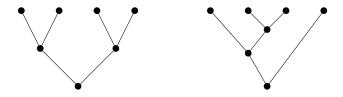


Now, flip the tree for R horizontally (so there is no crossover). Pair up corresponding leaves with an edge (representing that X_i in L is equal to X_i in R), and pair up the roots with an edge (as we want L = R).

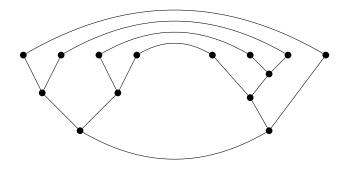




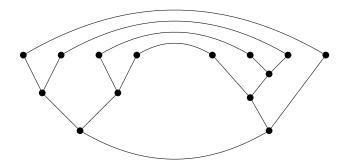
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By removing the leaf vertices, this forms a bridgeless cubic planar graph.

Suppose we have a sharp solution to L = R. We label the vertices with the result of the cross product immediately above it, ignoring signs.

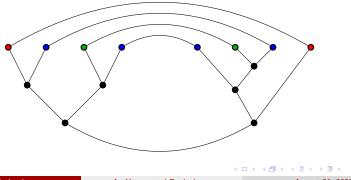
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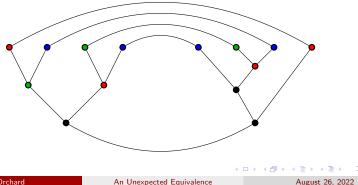
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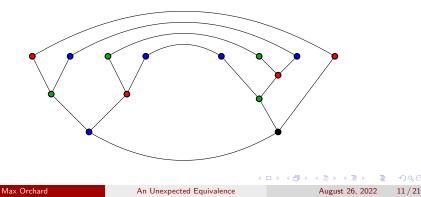


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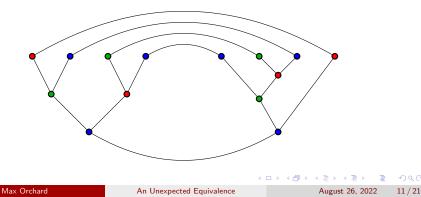
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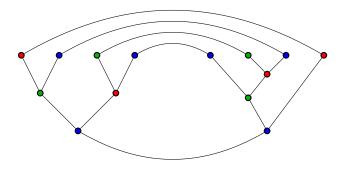


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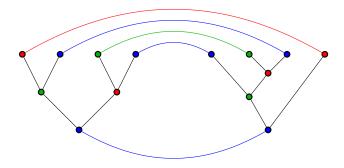


We can now obtain a Tait colouring, using the colour of the vertex at the "top" of the edge.



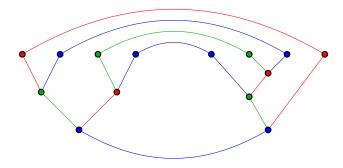
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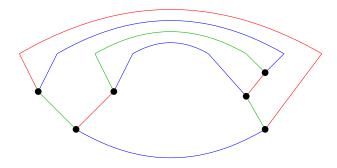
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Theorem

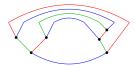
Given a sharp solution to L = R, we can obtain a Tait colouring of the associated bridgeless cubic planar graph.

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Proof.

Colour the graph as before.



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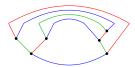
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Colour the graph as before. This is a 3-colouring as we have a sharp solution (so the only possible options for the vertices are $\{\pm i, \pm j, \pm k\}$) and we are ignoring signs.

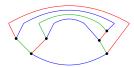


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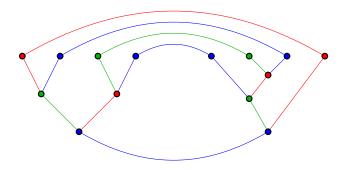
Colour the graph as before. This is a 3-colouring as we have a sharp solution (so the only possible options for the vertices are $\{\pm i, \pm j, \pm k\}$) and we are ignoring signs. It is a proper colouring due to the cyclic nature of the cross product on $\{i, j, k\}$ (ignoring signs).



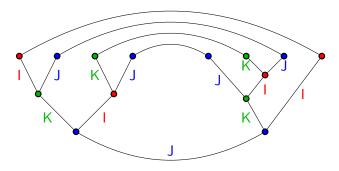
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Recall that the cross product is *anti-commutative* (i.e $a \times b = -(b \times a)$).

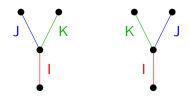
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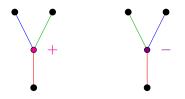
Recall that the cross product is *anti-commutative* (i.e $a \times b = -(b \times a)$). This means determining the sign is equivalent to determining the orientation of colours at a vertex.

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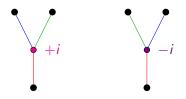


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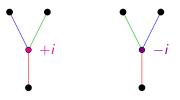
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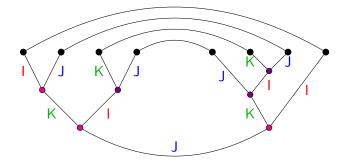
Because of this assignment, multiplying the labels for *L*'s tree will "give" the sign of $L(X_1, \ldots, X_n)$, and similarly for *R*'s tree. This follows from bilinearity.

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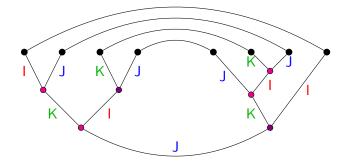
We now label the vertices of our graph using the orientation of I, J, K.

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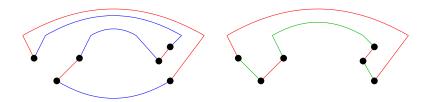
As the tree for R is flipped, we must flip the labelling on the right.

A formation is a graph formed from exactly two edge colours.

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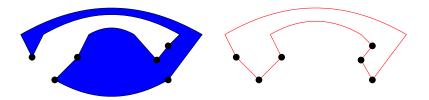
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A *formation* is a graph formed from exactly two edge colours. As the degree of each vertex in a formation is 2, a formation can be decomposed into a product of cycles.

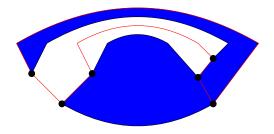


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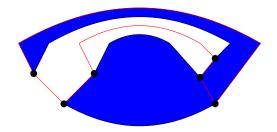


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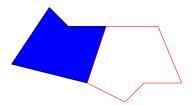


When two formations overlap, there are two different ways their edges can interact.

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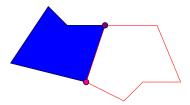
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A *bounce* occurs when both formations share an edge and the interior of each formation is entirely disjoint, or one is contained inside the other.

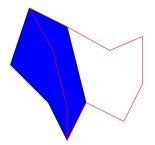


A *bounce* occurs when both formations share an edge and the interior of each formation is entirely disjoint, or one is contained inside the other.

Each bounce consists of a +i vertex and a -i vertex. Therefore, each bounce contributes 1 to the complex product of the labels.



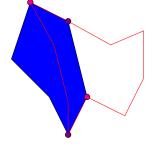
A crossing occurs when the interiors of both formations partially intersect.



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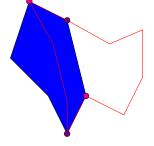
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Lemma

The complex product of the labels in a Tait colouring is 1.

Max Orchard

An Unexpected Equivalence

August 26, 2022 19 / 21

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Theorem

Given a Tait colouring of the associated bridgeless cubic planar graph, we can obtain a sharp solution to L = R.

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Let Z be the complex product for the tree T(L) of L, and W be the complex product for the tree T(R) of R.

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 $Z\overline{W} = 1 \implies Z = W.$

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$$Z = e(i^m), \quad W = e'(i^m).$$

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Thus e = e', and we have a sharp solution to L = R.

References

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- P. Rideout. *Cross Products and the Four Color Theorem*, 2020. Accessed from:

https://prideout.net/blog/kauffman/kauffman.pdf.

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